Veriﬁed Bidirectional Transformations by Construction

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Abstract. Bidirectional transformations (bx) are of increasing signiﬁcance in model-based engineering. Currently bx are deﬁned using a number of specialised transformation languages. In this paper we show how standard UML elements such as use cases and OCL constraints can be used to deﬁne bx, thus taking advantage of the wide industrial and educational adoption and support for UML. We deﬁne patterns and techniques for specifying bx so that they are correct by construction.

1 Introduction

Bidirectional model transformations (bx) provide a means to maintain a consistency relation $R$ between two models which may both change.

Bx are signiﬁcant for a wide range of model based engineering scenarios such as round-trip engineering; automatic consistency maintenance of multiple interrelated models, and application and database co-evolution [3].

Bx are characterised by a binary relation $R : SL \leftrightarrow TL$ between a source language (metamodel) $SL$ and a target language $TL$. $R(m, n)$ holds for a pair of models $m$ of $SL$ and $n$ of $TL$ when the models consist of data which corresponds under $R$. It should be possible to automatically derive from the deﬁnition of $R$ both forward and reverse transformations $R^\rightarrow : SL \times TL \rightarrow TL$ and $R^\leftarrow : SL \times TL \rightarrow SL$ which aim to establish/maintain $R$ between their ﬁrst (respectively second) and their result target (respectively source) models, given both existing source and target models.

Bx should satisfy two key conditions [16, 12]:

1. Correctness: the forward and reverse transformations derived from a relation $R$ do establish $R$: $R(m, R^\leftarrow (m, n))$ and $R(R^\rightarrow (m, n), n)$ for each $m : SL$, $n : TL$.

2. Hippocraticness: if source and target models already satisfy $R$ then the forward and reverse transformations do not modify the models: $R(m, n) \Rightarrow R^\rightarrow (m, n) = n$ and $R(m, n) \Rightarrow R^\leftarrow (m, n) = m$ for each $m : SL$, $n : TL$. 
Additionally, updates to a model by $R^+ \text{ or } R^-$ should be the minimal necessary to restore $R$: this is termed the principle of least change [13].

Section 2 describes current bx research, Section 3 describes some bx patterns, and Section 4 describes our approach to bx specification in UML. Section 5 provides a proof that the bx properties for $R^+$ are satisfied by transformations defined using our approach. Section 6 gives an evaluation of the approach.

2 Bx Specification Languages and Techniques

Bx are the subject of extensive theoretical research [16, 17, 4, 12] and an annual conference, BX. The most well-known transformation language which supports definition of bx is QVT-R [14]. However, QVT-R has had limited uptake because of its complex semantics and specification style, compared to simpler unidirectional MT languages such as ATL [6] and ETL [7]. Triple Graph Grammars (TGG) are a graph transformation-based formalism which supports bx definition, although with a restricted expression language compared to OCL. Other approaches, such as JTL [5], and extensions of TGG [1], use constraint-based programming techniques to interpret relations $P(s, t)$ between source and target elements as specifications in both forward and reverse directions. Model finding using Alloy is used to bidirectiona lise ATL in [12].

In this paper we describe an approach to bx definition using standard UML 2 elements (class diagrams, use cases, OCL) which avoids the use of specialised languages, and therefore is preferable from the viewpoint of industrial uptake, and for the long-term maintenance and reuse of bx. In our approach, unidirectional transformations $\tau$ are specified using use cases and OCL, then a statically-computed bx relation $Post_\tau \& Inv_\tau$ for $\tau$ is derived, together with a forward transformation $\tau^+$ extending $\tau$, and an inverse $\tau^-$ of $\tau^-$. The bx relation and transformations are derived from $\tau$ using a higher-order transformation. This is generally more efficient than the use of constraint programming and model finding.

3 Patterns for Bidirectional Transformations

We have identified a number of patterns for bx, which can be used to establish the bx correctness conditions [11]:

- **Auxiliary Correspondence Model**: maintain a detailed trace between source model and target model elements to facilitate change-propagation in source to target or target to source directions.

- **Cleanup before Construct**: for $R^-$, remove superfluous target model elements which are not in the transformation relation $R$ with any source elements, before constructing target elements related to source elements, and similarly for $R^+$.

- **Unique Instantiation**: Do not recreate elements $t$ in one model which already correspond to an element $s$ in the other model, instead modify the
data of \( t \) to enforce the transformation relation. Use key attributes to identify when elements should be created or updated.

All three of these patterns are used in our approach to define bx using UML. In addition, we have used three further patterns to structure bx:

- **Factor a Text-to-text bx by a Model-to-model bx**: this architectural pattern defines a text-to-text bx as a composition of parsing, model-to-model and model-to-text transformations in both directions.
- **Inverse Recursive Descent**: if a forward transformation \( \tau^{\rightarrow} \) is defined using structural recursion on the source language \( SL \), the reverse transformation \( \tau^{\leftarrow} \) may be derived from \( \tau^{\rightarrow} \) as a dual structural recursion on \( TL \).
- **Bx Data Correspondence**: a special case of Auxiliary Correspondence Model which defines an auxiliary association in \( SL \cup TL \) to express a bx relation concerning specific elements or basic values (e.g., a correspondence between expression operators in one language and the other).

4 Bidirectional Transformation Specification in UML-RSDS

UML-RSDS is a particular version of executable UML (xUML), based on UML 2 and OCL 2.4, with a formal semantics [8] and an established toolset [9]. Model transformations \( \tau \) are specified in UML-RSDS as UML use cases, defined declaratively by two main predicates, expressed in a subset of OCL:

1. Postconditions \( Post_{\tau} \) which define the intended effect of the transformation at its termination. These are an ordered conjunction \( C_1 \& \ldots \& C_n \) of OCL constraints (also termed rules in the following) and also serve to define a design (activity) and implementation of the transformation.
2. Invariants \( Inv_{\tau} \) which define expected invariant properties which should hold during the transformation execution. These can be derived from \( Post_{\tau} \), or can be specified explicitly by the developer. They are the constraints of the use case, considered as a UML classifier.

A UML activity \( stat(Cn) \) is derived by the UML-RSDS tools for each post-condition constraint \( Cn \). This activity is constructed so that it establishes \( Cn: [stat(Cn)]Cn \). The activity for \( Post_{\tau} \) is formed as a composition of the \( stat(Cn) \). In this paper, \( stat(Post_{\tau}) \) is the sequential composition of the \( stat(Cn) \).

Data-dependency relations between the \( Post \) constraints are important in ensuring the correctness of both unidirectional and bidirectional UML-RSDS transformations. Let \( rd(Cn) \) denote the read-frame of a constraint \( Cn \), the set of entity type names and data features that it reads, and \( wr(Cn) \) denote its write frame. These are defined in [8].

A dependency ordering \( Cn < Cm \) is defined between constraints by \( wr(Cn) \cap rd(Cm) \neq \{ \} \). “\( Cm \) depends on \( Cn \)”.

A use case with postconditions \( C_1, \ldots, C_n \) should satisfy the **syntactic non-interference** conditions:
1. If \( C_i < C_j \), with \( i \neq j \), then \( i < j \).
2. If \( i \neq j \) then \( \text{wr}(C_i) \cap \text{wr}(C_j) = {} \).

Together, these conditions ensure that the activities \( \text{stat}(C_j) \) of subsequent constraints \( C_j \) cannot invalidate earlier constraints \( C_i \), for \( i < j \).

A constraint is termed a type 1 constraint if its write and read frames are disjoint. Such constraints usually have an implementation as a bounded loop, as opposed to a fixed-point/unbounded iteration.

An invariant \( \text{Inv} \) and explicit inverse predicate \( \text{Post}^- \) can be derived mechanically from \( \text{Post} \) [9, 11]. This derivation is automated in the UML-RSDS tools. Tables 1 and 2 show some examples of inverses \( P^- \) of predicates \( P \). The transformation developer can also specify inverses for particular \( C_n \) by defining a suitable \( C_n^- \) constraint in \( \text{Inv} \), for example, to express that a predicate \( t.z = s.x + s.y \) should be inverted as \( s.x = t.z - s.y \).

<table>
<thead>
<tr>
<th>( P(s,t) )</th>
<th>( P^-(s,t) )</th>
<th>Condition</th>
</tr>
</thead>
<tbody>
<tr>
<td>( t.g = s.f )</td>
<td>( s.f = t.g )</td>
<td>Assignable features ( f, g )</td>
</tr>
<tr>
<td>( t.g = s.f \rightarrow \text{pow}(x) )</td>
<td>( s.f = t.g \rightarrow \text{pow}(1.0/x) )</td>
<td>( x ) non-zero, independent of ( s.f )</td>
</tr>
<tr>
<td>( t.b2 = \text{not}(s.b1) )</td>
<td>( s.b1 = \text{not}(t.b2) )</td>
<td>Boolean attributes ( b1, b2 )</td>
</tr>
<tr>
<td>( t.g = K \ast s.f + L )</td>
<td>( s.f = (t.g - L)/K )</td>
<td>( f, g ) numeric attributes</td>
</tr>
<tr>
<td>( K \neq 0 )</td>
<td></td>
<td></td>
</tr>
<tr>
<td>( R(s,t) \land Q(s,t) )</td>
<td>( R^-(s,t) \land Q^-(s,t) )</td>
<td></td>
</tr>
<tr>
<td>( s \rightarrow \forall x \mid P(x,y) )</td>
<td>( t \rightarrow \forall y \mid E \rightarrow \exists x \mid P^-(x,y) \land x : s) )</td>
<td>( E ) is the element type of ( s )</td>
</tr>
</tbody>
</table>

Table 1. Inverse of Predicates

In terms of the framework of [16], the source-target relation \( R_\tau \) associated with a UML-RSDS transformation \( \tau \) is \( \text{Post}_\tau \land \text{Inv}_\tau \). \( R_\tau \) is not necessarily bijective. The forward direction of \( \tau \) is normally computed as \( \text{stat}(\text{Post}_\tau) \): the UML activity derived from \( \text{Post}_\tau \) when interpreted procedurally [8]. However, in order to achieve the correctness and hippocratic properties, \( \text{Inv}_\tau \) must also be considered: before \( \text{stat}(\text{Post}_\tau) \) is applied to the source model \( m \) and target model \( n \), \( n \) must be cleared of elements which fail to satisfy \( \text{Inv}_\tau \). This is a case of the Cleanup before Construct pattern.

Given a typical postcondition \( C_i \) of the form

\[
S_i ::
SCond_i(self) \Rightarrow T_j \rightarrow \exists s \mid TCond_j(t) \land P_i(self, t))
\]

the corresponding invariant constraint is \( \text{Inv}_i \):

\[
T_j ::
TCond_j(self) \Rightarrow S_i \rightarrow \exists s \mid SCond_i(s) \land P_i(s, self))
\]
The general cleanup constraint derived from \( \text{Inv}_s \) is:

\[
T_j ::
\begin{align*}
\text{TCond}_j(\text{self}) & \land \\
\text{not}(\text{SCond}_i(\text{self}) & \land \text{P}_i(\text{self})) & \Rightarrow \text{self} \rightarrow \text{isDeleted}()
\end{align*}
\]

However, this constraint may delete more \( T_j \) instances than is necessary: not only instances which correspond (i) to deleted \( S_i \) instances, or (ii) to \( S_i \) instances that no longer satisfy \( S\text{Cond}_i \), but also (iii) to \( S_i \) instances satisfying \( S\text{Cond}_i \), but not \( P_i(\text{self}) \). In this last case, \( \text{self} \) should not be deleted, instead \( \text{stat}(P_i) \) should be performed to re-establish \( P_i \) between the modified \( s \) and \( \text{self} \). Therefore, some form of tracing is required, to distinguish these cases.

The Auxiliary Correspondence Model pattern can be used to introduce tracing relations between source and target entities. In general such relations can be many-many, and are implemented by an auxiliary association \( S_i \rightarrow s\text{trace} - \rightarrow t\text{trace} T_j \) linking corresponding elements in the two languages. In the transformation specification, \( s\text{.trace} \) can be referred to as \( s \rightarrow \text{equivalents}() \), and \( t\text{.trace} \) as \( t \rightarrow \text{equivalents}() \). The computational overhead of maintaining such relations can be substantial, and instead we usually apply the Auxiliary Correspondence Model pattern to introduce String-valued identity attributes (primary keys) for source and target entities to record and identify which source and target instances correspond: source instance \( s \) corresponds to target instance \( t \) when \( s\text{.sid} = t\text{.id} \). This also means that \( S_i(\text{t.id}) = s \) and \( T_j(s\text{.sid}) = t \). Finally, the Unique Instantiation pattern is used to modify existing objects in one model which correspond to a modified object in the other model, instead of creating a new object.

If identity attributes have been introduced using Auxiliary Correspondence Model, the cleanup constraints can be simplified to:

\[
T_j ::
\begin{align*}
\text{TCond}_j(\text{self}) & \land \\
\text{not}(\text{t.id} : S_i \rightarrow \text{collect(s.id)} & \Rightarrow \text{self} \rightarrow \text{isDeleted}())
\end{align*}
\]
which correspond to cases (i) and (ii) above. This version is also potentially more efficient than the original cleanup constraint. The second case is omitted if SCond_i is the default true. The updates of case (iii) are carried out in the Post constraints. \( \tau \rightarrow \) is defined to have postconditions consisting of the cleanup constraints, followed by Post, and hence it is an extension of \( \tau \); its postconditions imply those of \( \tau \).

The constraints of Post- need to be interpreted using Unique Instantiation in \( \tau \rightarrow \): the \( E \rightarrow \exists (e \mid P) \) quantifier in rule succedents should be interpreted as “create a new \( e : E \) and establish \( P \) for \( e \), unless there already exists an \( e : E \) satisfying \( P \)”. This is also known as ‘check before enforce’ semantics. If key attributes are defined, a constraint

\[
E :: \\
F \rightarrow \exists (f \mid f.fId = eId \& P)
\]

means that if \( eId : F \rightarrow \text{collect}(fId) \) (a corresponding \( F \) instance \( F[eId] \) already exists), then it is selected and modified to satisfy \( P \) – this performs case (iii) of the cleanup actions to re-establish Inv after a source model change. Only if \( eId \notin F \rightarrow \text{collect}(fId) \) is a new \( F \) instance created.

5 Correctness and Hippocraticness Properties for \( \tau \rightarrow \)

In this section we give a formal justification that bx defined using the UML specification approach described above do satisfy the bx properties.

We consider first the case of UML-RSDS transformations that satisfy the Structure Preservation pattern [10], in which source language entities \( S_i \) correspond to target language entities \( T_i \), and whose constraints are of type 1, satisfying syntactic non-interference. Identity attributes are used to implement Auxiliary Correspondence Model. This means that \( \tau \) is a separate-models transformation with source language \( S \) and target language \( T \), and postcondition Post is an ordered conjunction of constraints \( C_i \) each of the form:

\[
S_i :: \\
SCond_i(self) \Rightarrow \\
T_i \rightarrow \exists (t \mid t.tId = sId \& TCond_i(t) \& P_i(self, t))
\]

and Inv is a conjunction of constraints Inv_i of the form

\[
T_i :: \\
TCond_i(self) \Rightarrow \\
S_i \rightarrow \exists (s \mid tId = s.sId \& SCond_i(s) \& P_i(s, self))
\]
The simplified cleanup constraints for $\tau \rightarrow$ are therefore:

$$T_i ::
TCond_i(self) \& \text{not}(tId : S_i \rightarrow \text{collect}(sId)) \Rightarrow \text{self} \rightarrow \text{isDeleted}()$$

and

$$T_i ::
TCond_i(self) \& tId : S_i \rightarrow \text{collect}(sId) \& 
\text{not}(SCond_i(S_i[tId])) \Rightarrow \text{self} \rightarrow \text{isDeleted}()$$

The second constraint is omitted if $SCond_i$ is the default true predicate. We denote the collection of the cleanup constraints by $\text{Cleanup}$. Bx correctness holds, since the cleanup constraints together with $\text{Post}$ establish Inv: the Cleanup constraints remove any target model elements that fail to correspond by identity to source model elements in the domain of $\tau$, and $\text{Post}$ modifies target model elements that do correspond to valid source elements, in order to re-establish Inv. By construction, $\text{stat}(\text{Post})$ preserves Inv and establishes $\text{Post}$. Therefore, the sequential composition of $\text{stat}(\text{Cleanup})$ and $\text{stat}(\text{Post})$ establishes the bx relation $R$ as $\text{Post} \& \text{Inv}$.

Hippocraticness holds, since $\text{stat}(\text{Cleanup})$ has no effect if the models already satisfy $R$, and neither does $\text{stat}(\text{Post})$, due to the Unique Instantiation interpretation of the exists operator. The principle of least change is satisfied, since changes to the source model either lead to deletion of exactly those target elements (and their incident links) which cannot be modified to correspond to source elements (cases (i) and (ii) above), or to creation/modification of target elements, using $\text{stat}(C_i)$). But $\text{stat}(C_i)$ is designed as a minimally-intrusive update to the target model which will establish $C_i$. Interpreting a quantifier $T_j \rightarrow \text{exists}(t \mid P)$ using Unique Instantiation means that un-necessary creation of target elements is avoided.

A more complex case is Entity Merging, where two or more source entity types $S_i$, $S_j$ may map to the same target entity type $T_i$. Provided that the respective $TCond_i$ conditions are pairwise disjoint, the same construction and argument for the bx correctness apply as in the Structure Preservation case. Semantic non-interference of $\text{Post}$ holds, because the set of $T_i$ instances produced from $S_i$ is disjoint from the set produced from $S_j$, although syntactic non-interference fails. Entity Splitting involves one source entity type mapping to two or more different target entity types. There are two versions of this pattern: (i) two or more separate postconditions on $S_i$, with disjoint $SCond_i$ conditions, mapping to distinct target entity types; (ii) a single constraint which maps one $S_i$ instance to multiple linked instances of different $T_j$. The first case can be treated in a similar way to the Structure Preservation case, provided that syntactic or semantic non-interference of $\text{Post}$ holds. The second case is common in

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1 A use case satisfies semantic non-interference if for $i < j$: $C_i \Rightarrow \text{[stat}(C_j)]|C_i$, where $\text{[act]}P$ is the weakest-precondition of $P$ with respect to act. Syntactic non-interference implies semantic non-interference, but not conversely.
refinement transformations, and in some migrations. Constraints, eg., $C_1$, will have the form

$$S_1 ::
SCond_1(self) \Rightarrow T_1 \rightarrow exists(t1 \mid t1.t1Id = sId &
T_2 \rightarrow exists(t2 \mid t2.t2Id = sId & TCond_1(t1) &
TCond_2(t1, t2) & P_1(self, t1, t2)))$$

Both $t1$ and $t2$ are given id values $self.sId$. This should be the only constraint which creates either $T_1$ or $T_2$ instances (this restriction is implied by syntactic non-interference). $Inv_1$ is of the form

$$T_1 ::
TCond_1(self) & t2 : T_2 & t2.t2Id = t1Id & TCond_2(self, t2) \Rightarrow
S_1 \rightarrow exists(s \mid t1Id = s.sId & SCond_1(s) \ & P_1(s, self, t2))$$

The $t2 : T_2$ expression acts like an additional $T_2 \rightarrow forall(t2$ quantifier over the remainder of the constraint.

The cleanup constraints are:

$$T_1 ::
TCond_1(self) & t2 : T_2 & t2.t2Id = t1Id & TCond_2(self, t2) &
not(S_1 \rightarrow exists(s \mid t1Id = s.sId & SCond_1(s))) \Rightarrow
self \rightarrow isDeleted() & t2 \rightarrow isDeleted()$$

The effect of such a bx is that a ‘cluster’ of target instances are related to each source instance, with all elements of the cluster sharing the source instance identity value.

Again, $stat(\text{Post})$ establishes $\text{Post}$ and preserves $Inv$, whilst $stat(\text{Cleanup})$ together with $stat(\text{Post})$ establishes $Inv$. Thus correctness follows. Hippocraticness holds because target model pairs (or clusters) $t1, t2, ...$ are only deleted if they do not correspond to a source instance. Unique Instantiation can apply to multiple exists quantifiers in the same manner as to single quantifiers. The principle of least change follows from the minimality of $stat(C_i)$ as an activity to establish $C_i$.

Finally, the Map Objects before Links pattern can be used. This separates the construction of target instances, and the linking of these instances. The linking constraints have the form:

$$S_1 ::
T_j[idS_1].rr = \text{TRef}[r.idSRef]$$

These define target model association ends $rr$ from source model association ends $r$, looking-up target model elements $T_j[idS_1]$ and $\text{TRef}[r.idSRef]$ which have already been created by a preceding constraint. However, the linking constraints are not invariants of the transformation, and are omitted from $Inv$. The cleanup constraints for such $r$ are defined as for the preceding cases, and by a similar
argument, bx correctness, hippocraticness and least change can be shown. The
re-establishment of Post and Inv for corresponding source and target elements
is now performed by possibly a number of separate Post constraints.

A similar analysis applies to the reverse transformation $\tau^\rightarrow$. This has post-
condition $Post^\rightarrow$, the explicit form of Inv. $Post^\rightarrow$ may involve view updates [2]
of the source language SL, based on changes to TL elements. For example, predic-
tates $s.f\rightarrow last() = t.g$ or $s.f\rightarrow select(P_1) = t.g$.

$stat(Post^\rightarrow)$ establishes Inv, whilst the cleanup constraints together with
$Post^\rightleftarrows$ establish Post. Thus $\tau^\rightarrow$ satisfies correctness. Hippocraticness follows
since neither the cleanup or $Post^\rightarrow$ constraints modify the source model if $R$
already holds.

6 Evaluation

The approach described in this paper enables transformation developers to fo-
cus on writing correct postconditions $Post_\tau$ for a conventional (unidirectional)
forward transformation $\tau$. From these, the necessary invariants and cleanup con-
straints for bx execution mode can be automatically generated by a higher-order
transformation, producing a bx relation $Post_\tau \& Inv_\tau$ and forward and reverse
transformations with respect to this relation. The forward transformation ex-
tends $\tau$. This is the same approach as taken for the bidirectionalisation of ATL
in [12]: from a unidirectional forward ATL transformation $\tau$, a bx relation $T$
and transformations $T^\leftarrow$ and $T^\rightarrow$ are derived, where $T^\rightarrow$ extends $\tau$. In contrast
to [12] we use the minimality of the activity $stat(P)$ to achieve the principle
of least change, rather than reliance upon an external formalism (Alloy). Meta-
model constraints are incorporated into the bx relation by [12]. Our approach is
instead to prove separately that $Post_\tau \& Inv_\tau \Rightarrow Th_{TL}$, where $Th_{TL}$ are the
target language constraints.

By using standard UML notations and concepts, our approach should en-
able bx specifications to be retained and maintained independently of particular
transformation technologies. The use of OCL for transformation specification fa-
cilitates the automated derivation of transformation invariants and reverse rules.
Our approach scales up to transformations of practical size, and we have applied
it to a complex UML to C code generator with over 150 mapping rules. The ap-
proach defined here could also potentially be used with other similar specification
formalisms, such as Alf/UML [15] and EVL (http://www.eclipse.org/epsilon/doc/evl).

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